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# Time Slot based Heavy Changer Prediction on P4 Switches

P4交换机上基于时隙的Heavy Changer预测研究

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# Contents

01 Backgrounds, Motivation and Challenges

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02 System Design and Implementation

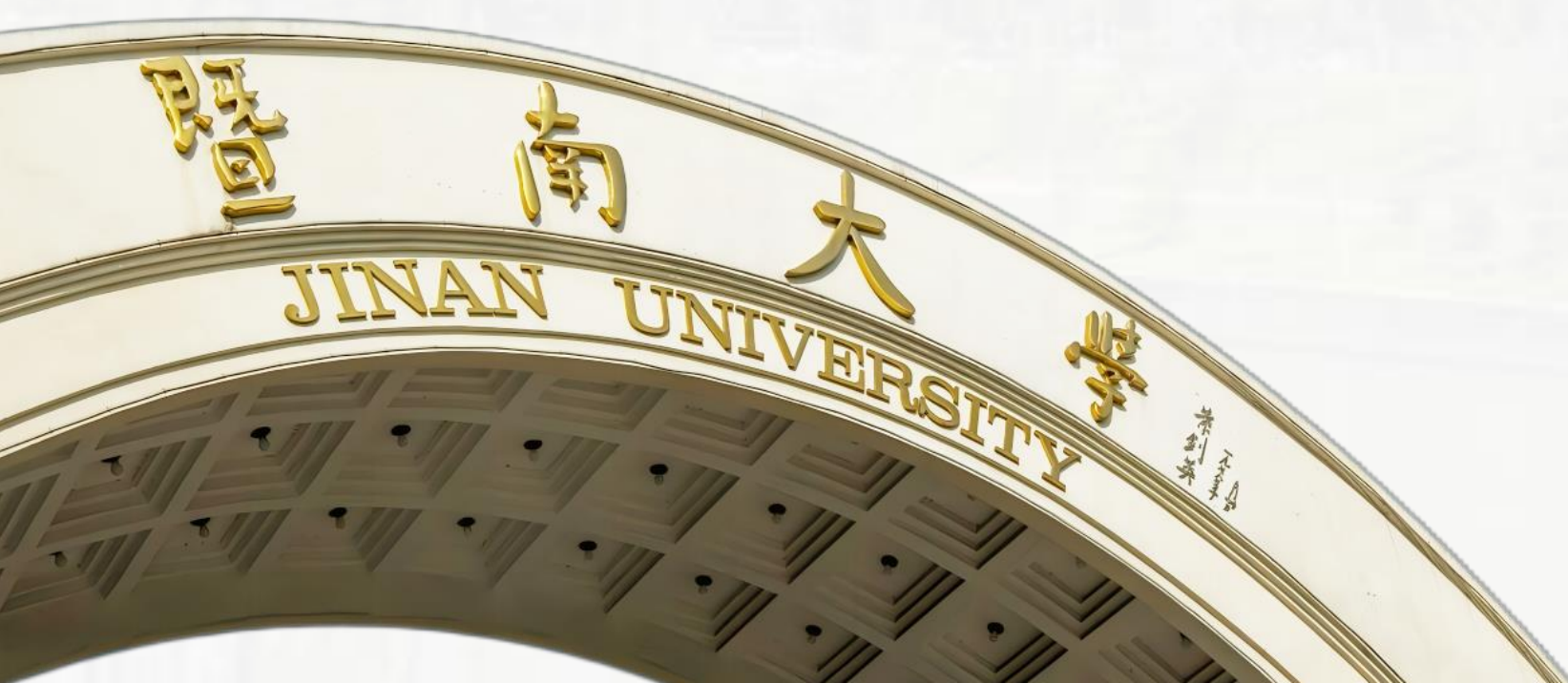
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03 Evaluation

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04 Conclusion

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# Backgrounds

## Heavy Changer, Time Slot, Ahead Definition

- ***Slot\_size***: time slot duration (in seconds)
- ***Ahead***: future slot count
- *time slots* or *epochs* or *slots*: regular time intervals
- $(x, v_x)$   $x$  flow identifier,  $v_x$  packet size
- $\phi (0 < \phi < 1)$  predefined heavy changer threshold
- $D$  absolute change in total flow  $D(x)$  absolute change in packets size
- Heavy changer:  $D(x) > \phi D$
- Heavy changers signal congestion or malicious attacks.



# Motivation and Challenges

## Motivation

- **Real-time** identification of heavy changers

Detection methods need packets

- **Avoids extra** control-data plane **communication**

Method Implement on the control plane need extra communication

## Challenges

- **Skewed Data Distribution:**

Imbalance data problem

Overrepresentation of normal flows

- **Data Plane Constraints:**

Limited memory and processing power

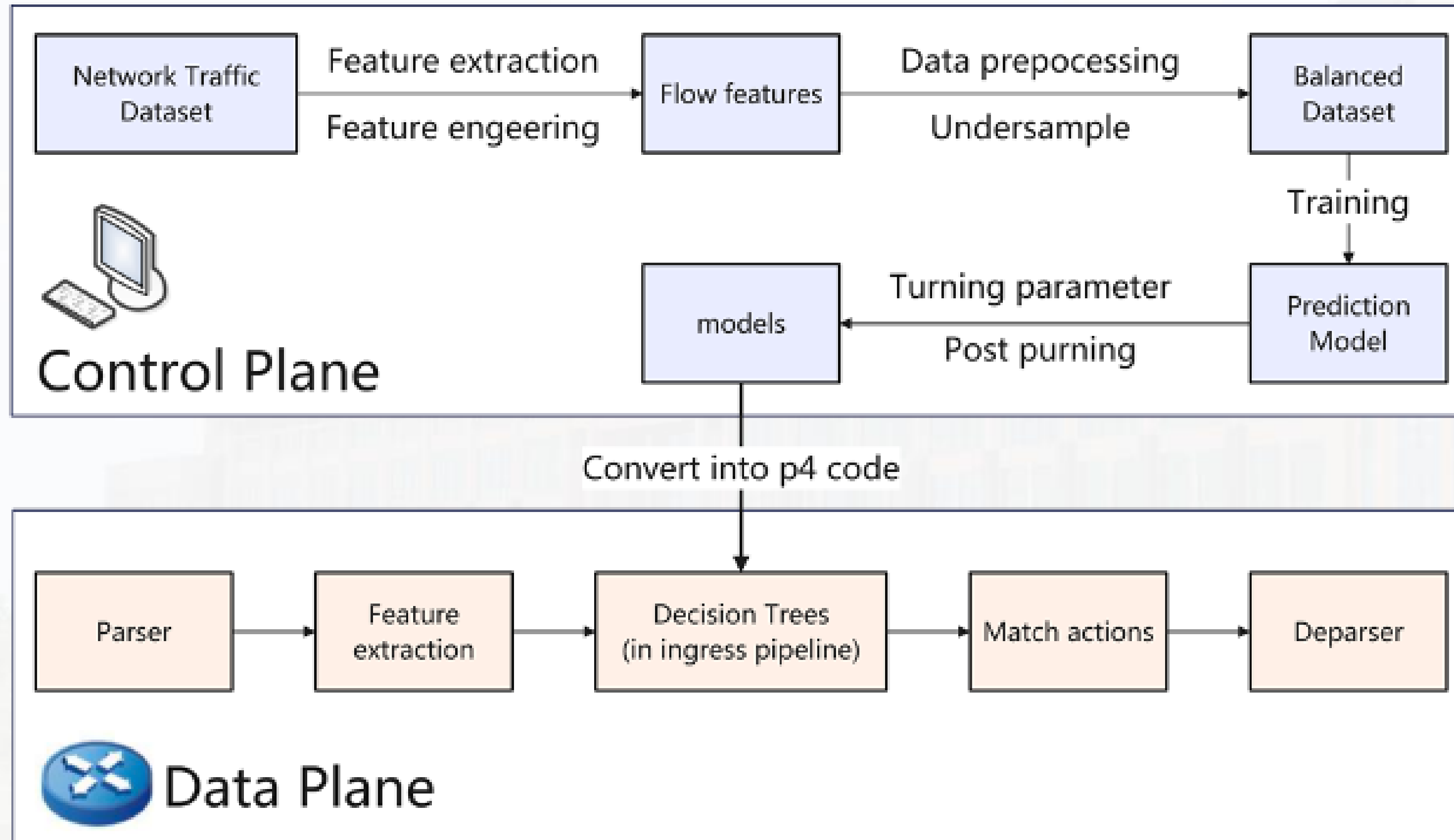
No floating-point operations, loops, division

- **P4 Language Limitations:**

Access only to packet metadata

Cannot predict across multiple flows

# System Overview



## Two Stages

- Training in the control plane

Mitigate data imbalanced problem

- Implement in the data plane

Convert not allowed operation

Using register to manage slots, timestamp

**Figure 1.** pChanger Procedure

# Training Process

## Feature

- Two Types of Features:

**Packet-level:** Dport, Sport, Proto

Extracted from metadata in P4 parse

**Flow statistics:** Length\_max, Length\_min, Length\_mean

IAT\_max, IAT\_min, IAT\_mean

Require register

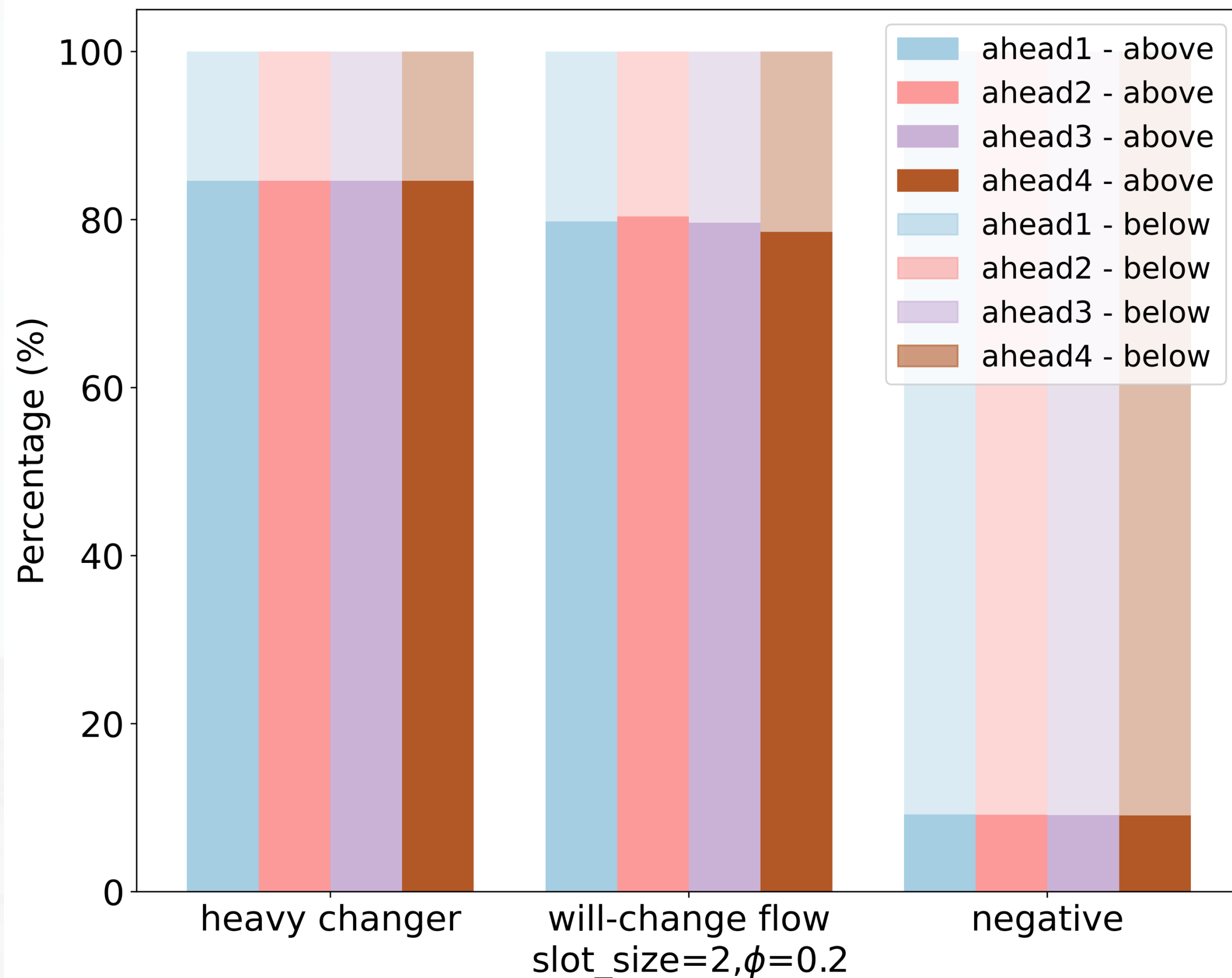
ACK\_flag, FIN\_flag, SYN\_flag

PSH\_flag, RST\_flag, ECE\_flag

Packet\_count

- **Removed Identifiers:** Excluding destination IPs, and source IPs.
- **Eliminated Timestamps:** Ensure heavy flow detection is time-independent.

# Data Filtering Strategies



- **Strategy 1:** Eliminate flows with fewer than 10 packets.
- **Strategy 2:** Eliminate flows with fewer than 20 packets.
- **Result:** Higher percentage of significant flows (heavy changers and will-change flows) retained.

**Figure 2.** Strategy 2 example: slot\_size 2, threshold 0.2



# Methods to mitigate imbalanced data problem

## •One-sided Selection

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**Algorithm** One-sided Selection↵

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**Input:**↵

$S$ : the initial training set.↵

$C$ : a stable subset composed of every positive from  $S$  alongside a negative chosen randomly↵

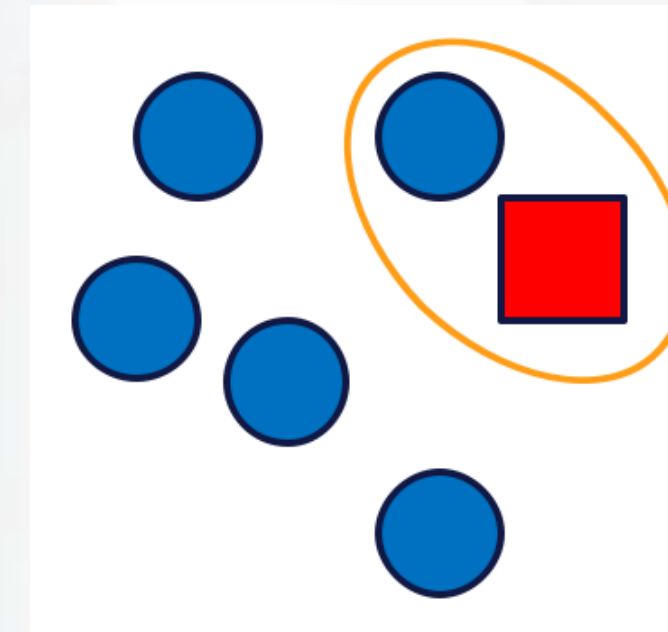
**Output:**↵

$T$ : reduced cleaner training subset↵

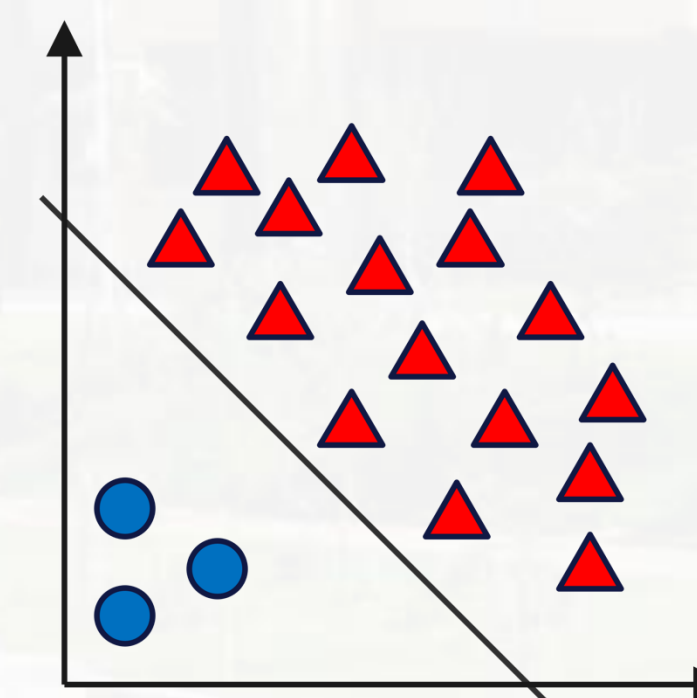
1. **For** sample  $s$  in  $S$  **do**↵
  2.     Classify  $s$  with  $C$  using 1-NN and verify label accuracy↵
  3.     **If**  $s$  is misclassified, move  $s$  into  $C$ ↵
  4. Eliminate all negative examples in  $C$  involved in Tomek links↵
  5.  $T \leftarrow C$ ↵
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Tomek Link to identify borderline and noisy examples (either marginal or noisy)

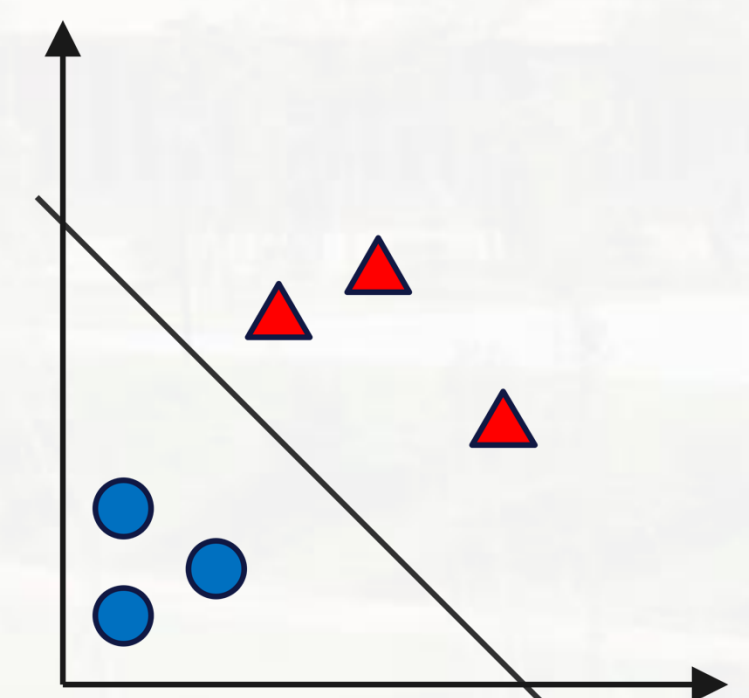
Two samples with different labels and no third sample are closer to them than they are to each other.



## • Under Sampling



Imbalanced data

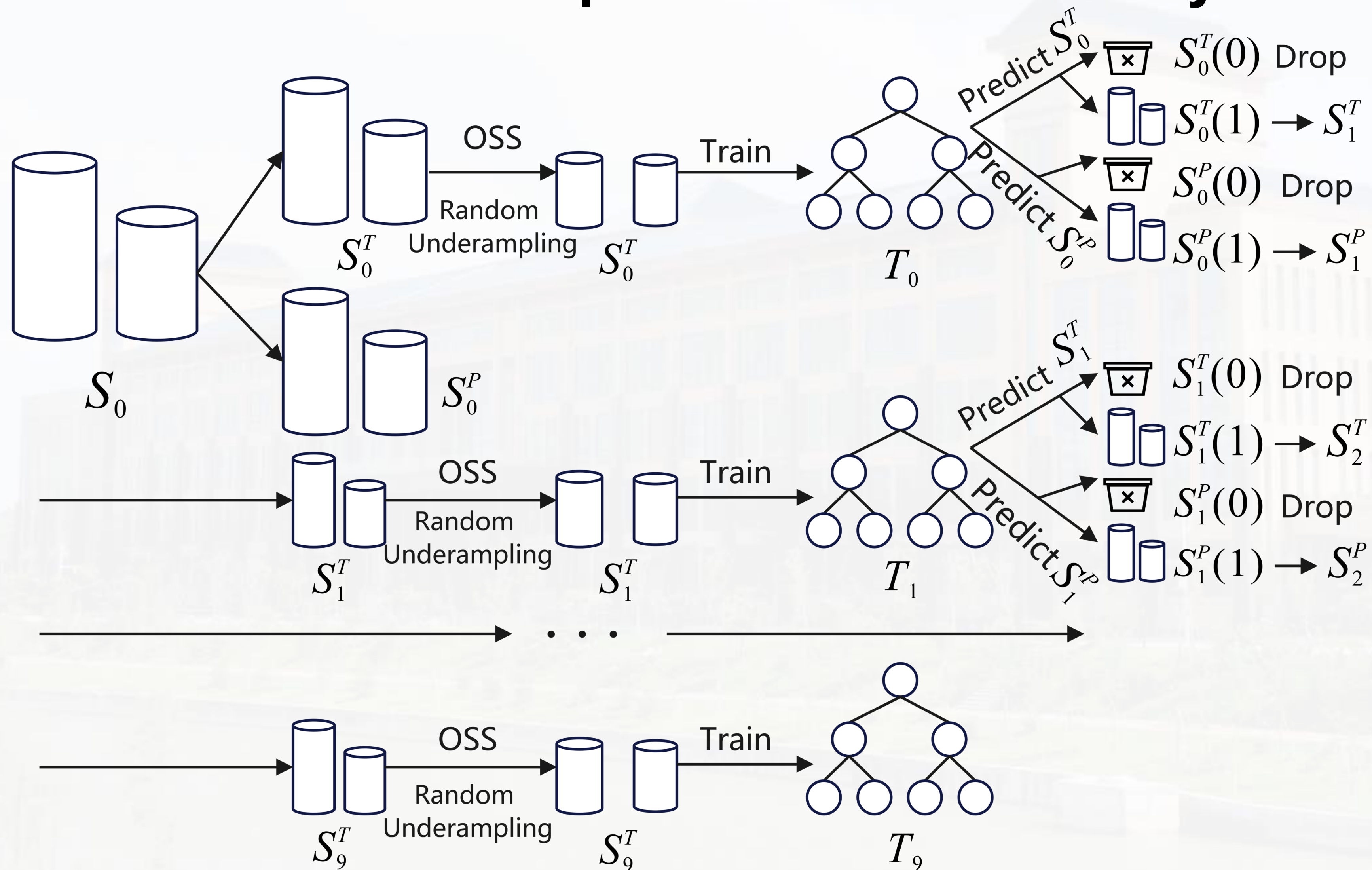


Undersampling



# Multi-stage model training process Post pruning

**Aim: Enhance subsequent models' accuracy**



# Solution of Data Plane Constraints Challenges

## Eliminate floating-point operations and division

- **Eliminate floating-point:**

Round down to the nearest integer

- **Eliminate division:**

Use exponentially weighted moving average (EWMA) and Set  $\alpha=0.5$

Formula becomes  $S_t = \frac{1}{2}y_t + \frac{1}{2}S_{t-1}$  i.e.,  $S_t = (y_t + S_{t-1}) >> 1$

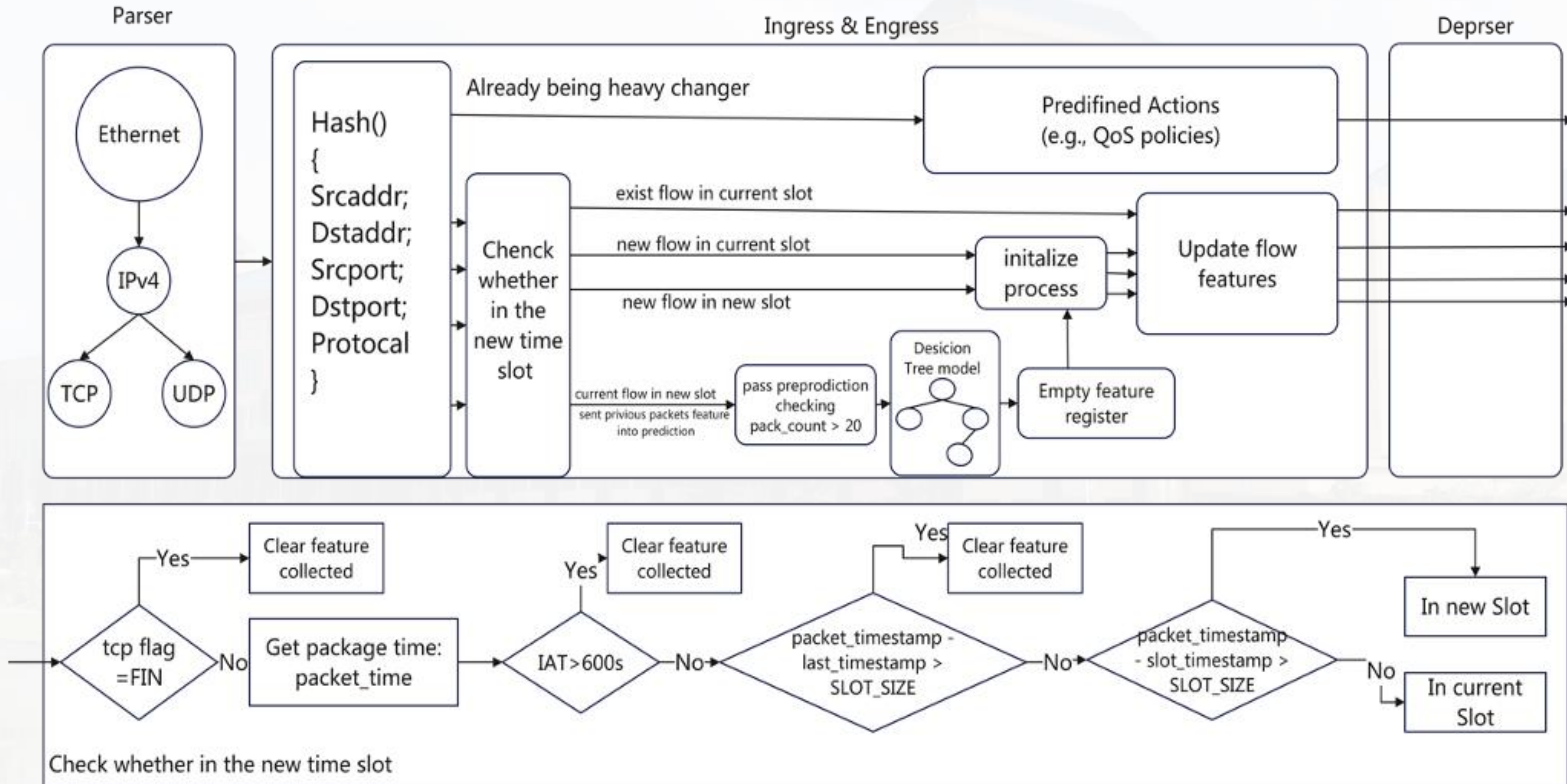
when  $t > 1$

- 
- The diagram illustrates the execution of four flows (A, B, C, D) across three time slots (Slot 1, Slot 2, Slot 3). The flows are represented by colored rectangles indicating their duration:
- Flow A (Blue):** Executes in Slot 1 and Slot 3.
  - Flow B (Teal):** Executes in Slot 2.
  - Flow C (Yellow):** Executes in Slot 1 and Slot 3.
  - Flow D (Red):** Executes in Slot 3.
- Vertical dashed lines separate the slots, and a horizontal axis at the bottom marks the boundaries of Slot 1, Slot 2, and Slot 3.



# Implementation in the data plane

## pChanger Pipeline Overview





# 4. Evaluation

## 4.1 Experiment Setup

### 1) Testbed:

- Server: Virtual server with 4 cores Intel i7-8550U 1.80GHz CPU, 20GB RAM.
- OS: Ubuntu 20.04.6 LTS.
- Datasets: Loaded into memory to avoid I/O overhead.

For P4 evaluation

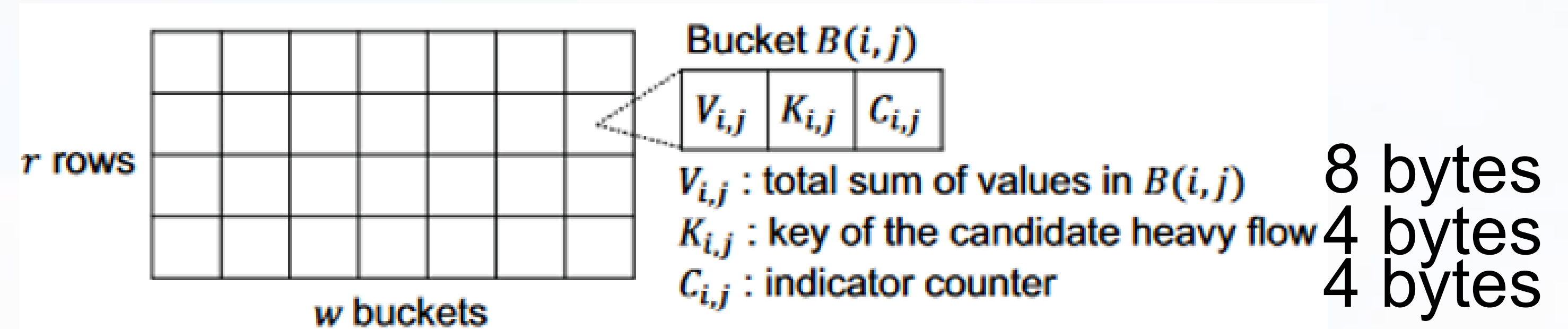
- Switch: bmv2 simple switch, input: veth2, output: veth4.

### 2) Dataset:

- Source: UNIV1 dataset from IMC 2010.
- Preprocessing: Payload nulled, IP anonymized with SHA1
- Division: Time slots of 2 seconds

For P4 evaluation

- Experiment Data: First 6 slots with slot\_size=2s, total 29,393 packets.



### 3) Methodology:

- Ground Truth: Mark true heavy changers based on threshold.
- Comparison: MV-Sketch.
- Memory Limiting: Adjust sketch width, use resource.setrlimit for pChanger.
- Feature Extraction: EWMA for averages counting

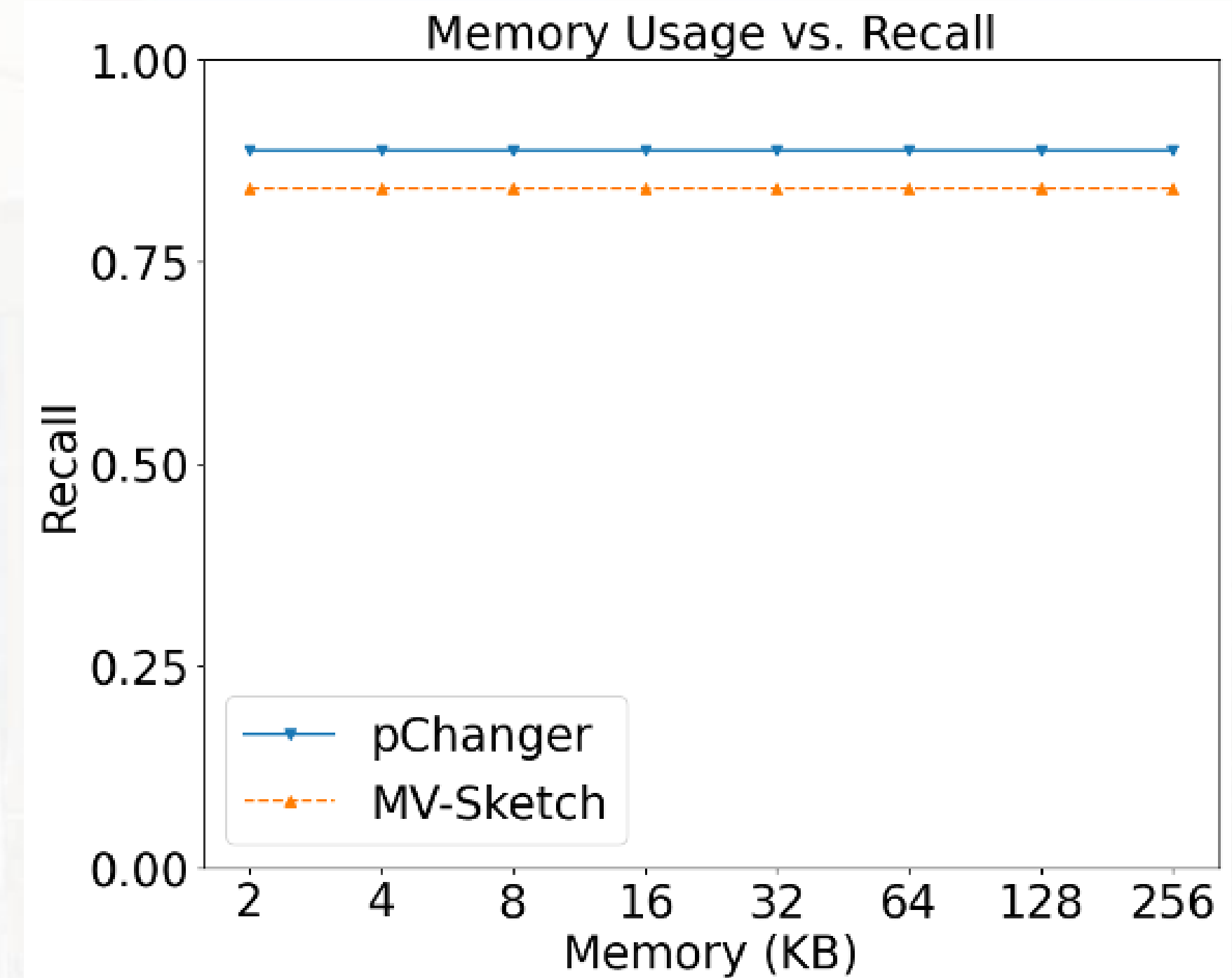
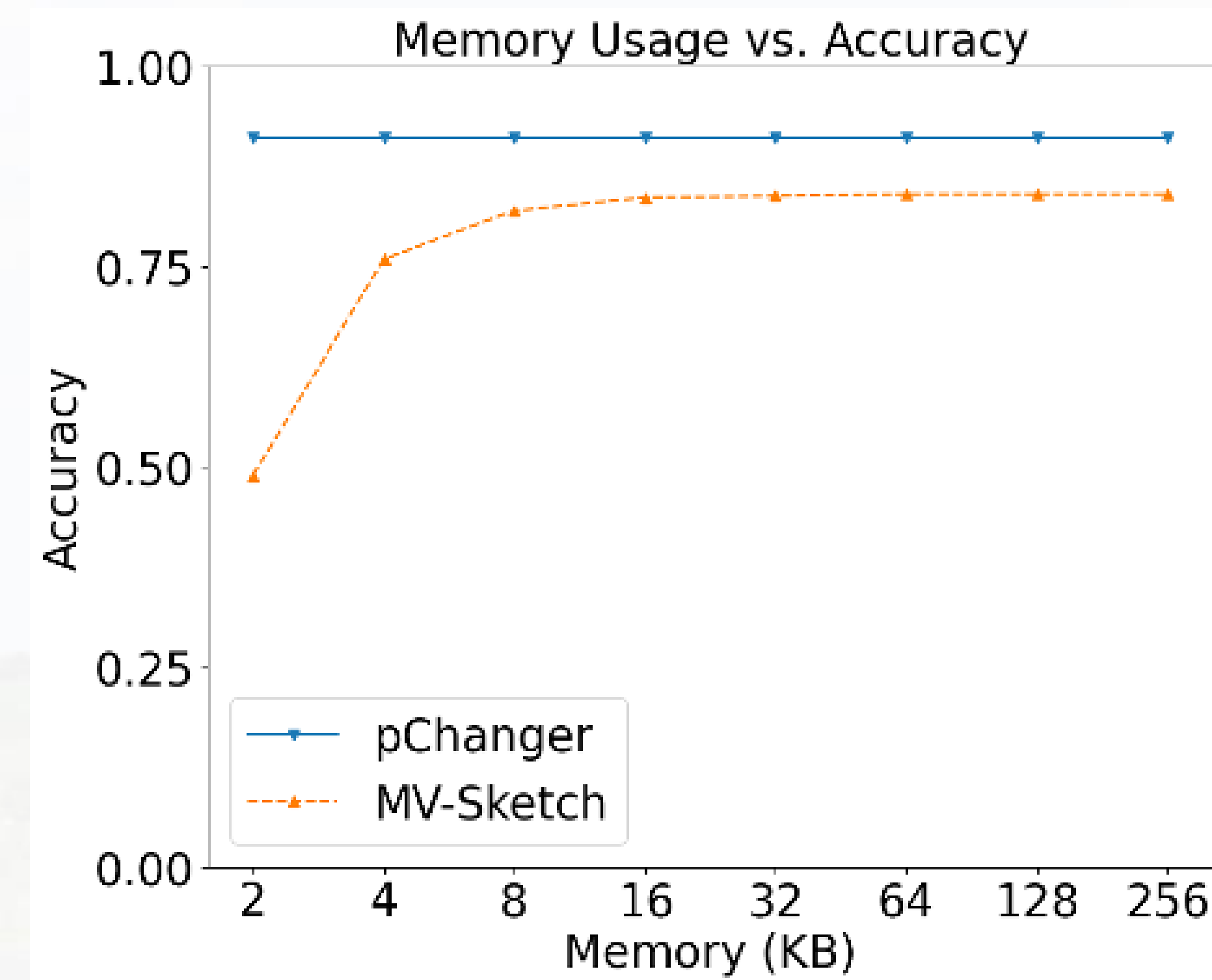
### 4) Metrics:

- Accuracy: Correctly identified flows.
- Precision: True heavy changers among reported flows.
- Recall: True heavy changers reported.
- F1-score: Harmonic mean of precision and recall.
- Update Throughput: Packets processed per second (Mbits/s).

# Results and analysis

## Experiment based on traces

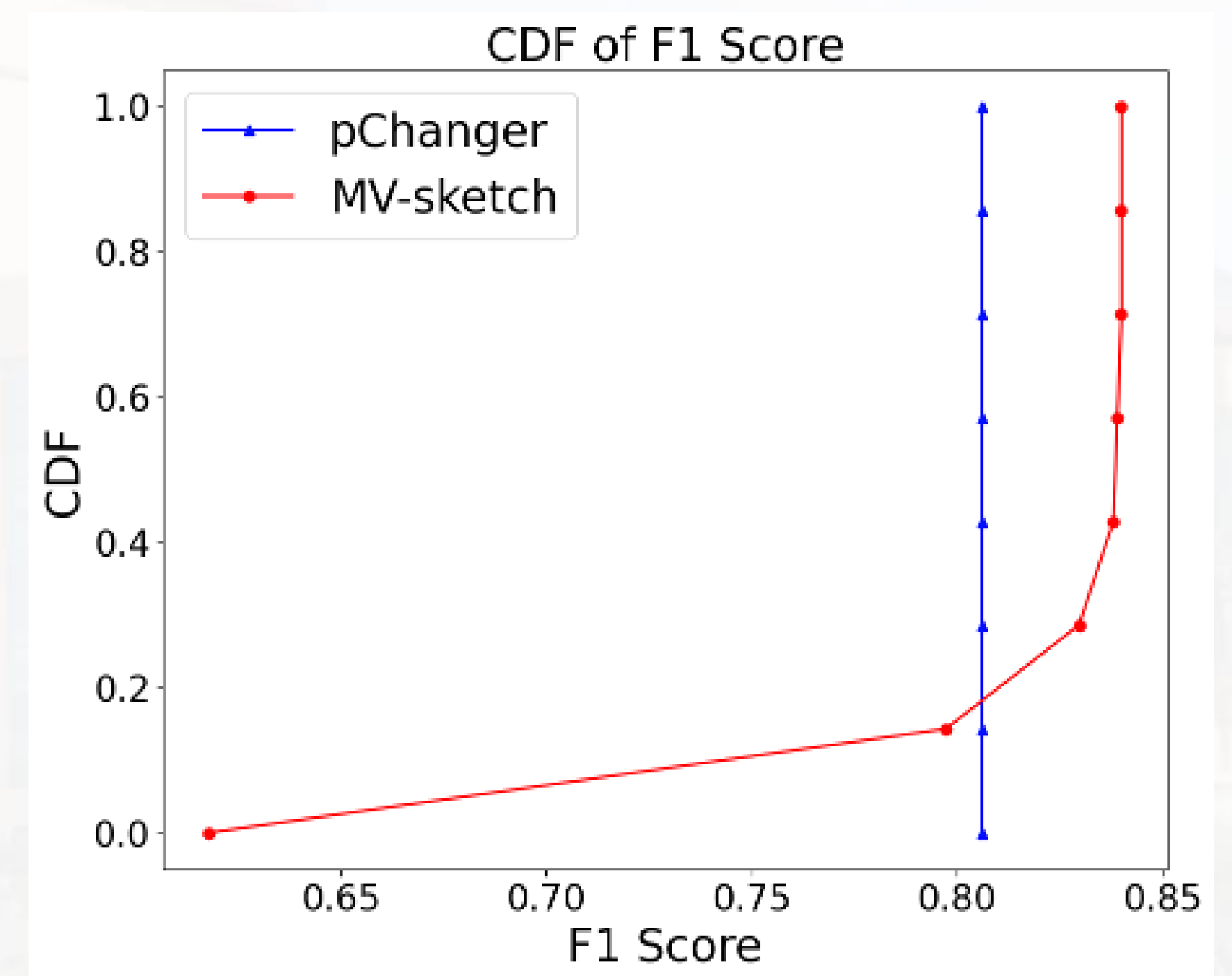
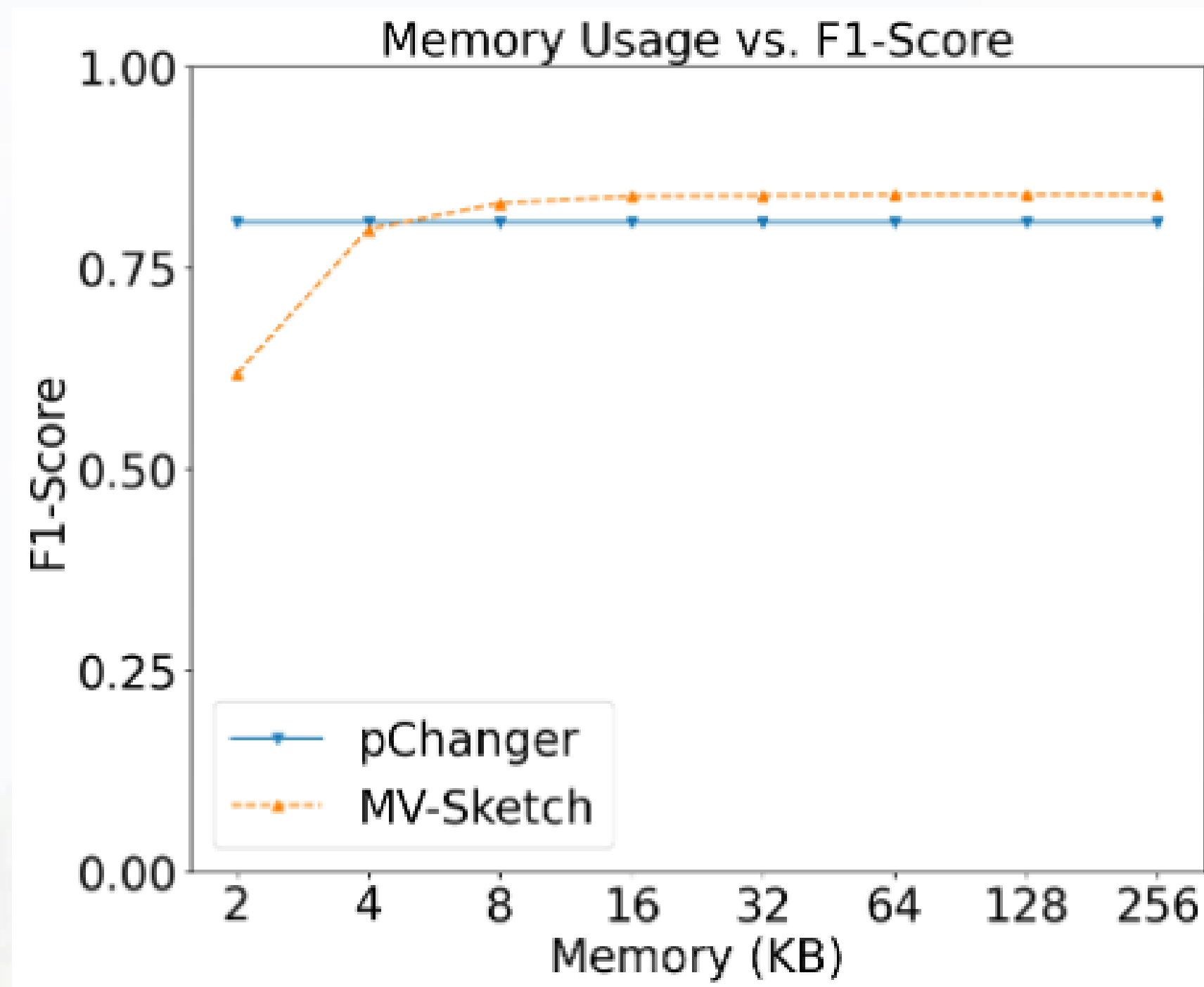
pChanger consistently maintains a high recall of 0.89 across all memory sizes, outperforming MV-Sketch, which improves from 0.81 to 0.84 as memory increases.



pChanger maintains high accuracy across all memory capacities, outperforming MV-Sketch, which improved accuracy only up to 4KB

# Results and analysis

## Experiment based on traces



pChanger outperforms MV-Sketch at low memory, but MV-Sketch surpasses pChanger as memory increases.

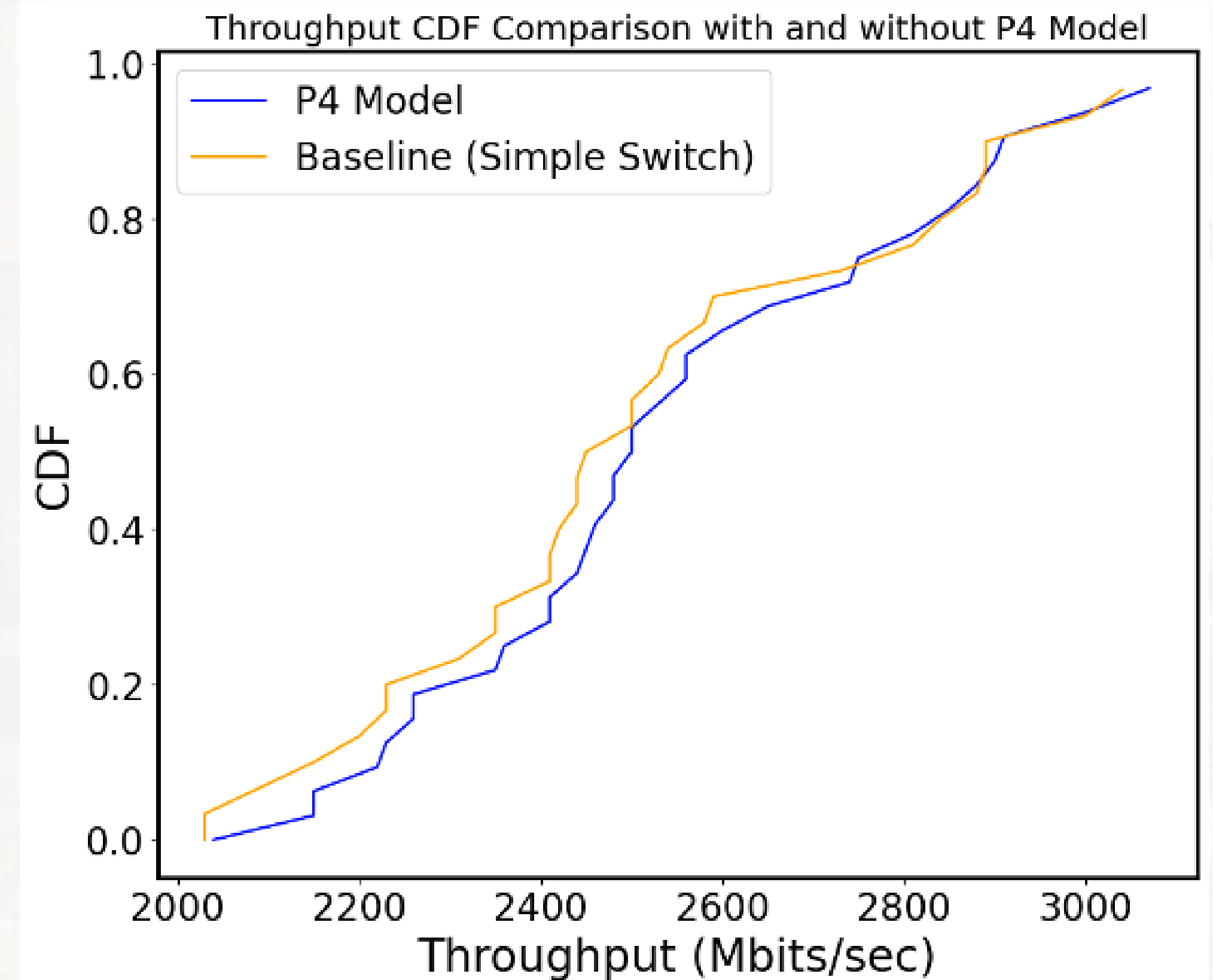
# Results and analysis

## Experiment on P4 Simple Switch

pChanger can run prediction at line rate

| Method   | Accuracy | Recall | Precision | F1-Score |
|----------|----------|--------|-----------|----------|
| pChanger | 0.9997   | 0.7500 | 0.6000    | 0.6667   |

- **Findings:**
  - High accuracy due to a small number of packets.
  - Rounding floating-point numbers reduces performance.
  - Predicts up to 4 slots in advance





# 5. Conclusion and Future works

## pChanger: programmable heavy changer prediction based on time slots

- **Advantages:**

- Reduces network congestion by predicting before heavy changers occur.
- Timeliness: Processes data in time slots.
- High accuracy and recall.
- Low memory usage doesn't impact throughput.
- Enables linear prediction.

- **Limitations:**

- Can't utilize extra memory effectively.
- Generalization issues exist

- **Future Work:**

- Explore multiple decision trees or random forests.
- Introduce more trees when extra memory is available.
- Allow controllers to dynamically configure and modify ML models.
- Combine ML models with sketch-based methods.